# A Theory of Spherical Diagrams CCCG 2022 

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Toronto - August 27, 2022

## 3D Art Gallery Problem

Given a polyhedron in $\mathbb{R}^{3}$, choose a (preferably small) set of vertices or edges that collectively see its whole interior.


These are called vertex guards and edge guards.

## Vertex-guarding polyhedra

The Art Gallery Problem for vertex guards may be unsolvable, even in some orthogonal polyhedra:


Some points in the central region are invisible to all vertices!

## Point-guarding polyhedra

Even if guards are not constrained to lie on vertices, there are (orthogonal) polyhedra that require $\Omega(n \sqrt{n})$ point guards!

outer view

cross section

## Edge-guarding polyhedra

These observations justify the study of edge guards.


Problem 1. How many edge guards are needed for a polyhedron?
Problem 2. Assuming that there is a point guard on every vertex of a polyhedron, how many additional edge guards are needed?

## Spherical Occlusion Diagrams: Introduction



When polygons in $\mathbb{R}^{3}$ are orthographically projected onto a sphere, their edges become arcs of great circles.

## Spherical Occlusion Diagrams: Introduction



Moreover, when a polygon is partially hidden (i.e., "occluded") by another, in the projection there are arcs feeding into other arcs.

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## Spherical Occlusion Diagrams: Introduction



If in an arrangement of polygons all vertices are occluded, then their edges project into a "Spherical Occlusion Diagram".

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## Spherical Occlusion Diagrams: Introduction



In particular, this applies to polyhedra: if all vertices are occluded, then the 1-skeleton projects into a Spherical Occlusion Diagram.

## Spherical Occlusion Diagrams: Definition



A Spherical Occlusion Diagram, or just "Diagram", is a finite non-empty collection of arcs of great circles on the unit sphere.

## Spherical Occlusion Diagrams: Definition



All arcs in a Diagram must be internally disjoint.

## Spherical Occlusion Diagrams: Definition



Both endpoints of each arc in a Diagram must lie in the interiors of some other arcs in the Diagram (every arc "feeds into" two arcs).

## Spherical Occlusion Diagrams: Definition



All the arcs in a Diagram that feed into the same arc must reach it from the same side.

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## Spherical Occlusion Diagrams: Examples

## Diagram axioms:

1. Arcs are internally disjoint.
2. Each arc feeds into two arcs.
3. All arcs that feed into the same arc
 reach it from the same side.

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## Spherical Occlusion Diagrams: Basic properties

## Proposition

No arc in a Diagram is longer than a great semicircle.


Proof. Otherwise it would have arcs feeding into it from both sides.

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## Spherical Occlusion Diagrams: Basic properties

## Corollary

No two arcs in a Diagram feed into each other.


Proof. Otherwise they would be longer than a great semicircle.

## Spherical Occlusion Diagrams: Basic properties

## Proposition

A Diagram partitions the sphere into convex regions (or "tiles").


Proof. Two points in the same region can be connected by a chain of arcs of great circles that does not intersect the Diagram.

## Spherical Occlusion Diagrams: Basic properties

## Proposition

A Diagram partitions the sphere into convex regions (or "tiles").


The arc joining the first and the third vertex of the chain do not intersect the Diagram, either...

## Spherical Occlusion Diagrams: Basic properties

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A Diagram partitions the sphere into convex regions (or "tiles").

...Otherwise, following the Diagram we would intersect the first two arcs in the chain, which is impossible by assumption.

## Spherical Occlusion Diagrams: Basic properties

## Proposition

A Diagram partitions the sphere into convex regions (or "tiles").


So we can simplify the chain, reducing it by one arc. Inductively repeating this reasoning, we can reduce the chain to a single arc.

## Spherical Occlusion Diagrams: Basic properties

## Proposition

A Diagram partitions the sphere into convex regions (or "tiles").


Any two points in the region are connected by an arc of a great circle that does not intersect the Diagram; hence, it is convex.

## Spherical Occlusion Diagrams: Basic properties

## Corollary

Every Diagram is connected.


Proof. If there are two connected components, each of them is a Diagram. So, one is contained in a tile $\mathcal{F}$ determined by the other.

## Spherical Occlusion Diagrams: Basic properties

## Corollary

Every Diagram is connected.


Take an arc in $\mathcal{F}$ with endpoints close to the first component that intersects the second component.

## Spherical Occlusion Diagrams: Basic properties

## Corollary

Every Diagram is connected.


The arc can be replaced by a chain that intersects neither connected component of the Diagram.

## Spherical Occlusion Diagrams: Basic properties

## Corollary

Every Diagram is connected.


So its endpoints are in the same tile determined by the whole Diagram, and this tile cannot be convex.

## Spherical Occlusion Diagrams: Basic properties

## Proposition

A Diagram with $n$ arcs partitions the sphere into $n+2$ tiles.


Proof. A Diagram induces a planar graph with $v$ vertices and $n+v$ edges. By Euler's formula, $f+v=n+v+2$, hence $f=n+2$.

## Spherical Occlusion Diagrams: Construction



How can we automatically generate large classes of Diagrams?

## Spherical Occlusion Diagrams: Construction



Start from a subdivision of the sphere into strictly convex tiles, where each tile has an even number of edges.

## Spherical Occlusion Diagrams: Construction



Note that the 1-skeleton of the tiling is bipartite, because it has no odd cycles.

## Spherical Occlusion Diagrams: Construction



We can turn each vertex of the tiling into a "swirl" going clockwise or counterclockwise according to the bipartition of the 1 -skeleton.

## Spherical Occlusion Diagrams: Construction



This operation defines a natural correspondence between even-sided spherical tilings and so-called swirling Diagrams.

## Swirling Diagrams: Examples

This method enables the automatic construction of swirling Diagrams from convex tilings of the sphere or convex polyhedra.


Prisms with even-sided bases

## Swirling Diagrams: Examples

This method enables the automatic construction of swirling Diagrams from convex tilings of the sphere or convex polyhedra.


Truncated antiprisms

## Swirling Diagrams: Examples

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Truncated bipyramids with even-degree vertices

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Trapezohedra

## Swirling Diagrams: Examples

This method enables the automatic construction of swirling Diagrams from convex tilings of the sphere or convex polyhedra.


Rhombic dodecahedron

## Swirling Diagrams: Examples

This method enables the automatic construction of swirling Diagrams from convex tilings of the sphere or convex polyhedra.


Deltoidal icositetrahedron

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This method enables the automatic construction of swirling Diagrams from convex tilings of the sphere or convex polyhedra.


Rhombic triancontahedron

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Deltoidal hexecontahedron

## Swirling Diagrams: Examples

This method enables the automatic construction of swirling Diagrams from convex tilings of the sphere or convex polyhedra.


Truncated cuboctahedron

## Swirling Diagrams: Examples

This method enables the automatic construction of swirling Diagrams from convex tilings of the sphere or convex polyhedra.


Truncated icosidodecahedron

## Swirling Diagrams: Alternative definitions



As we saw, a swirl in a Diagram is a cycle of arcs such that each arc feeds into the next going clockwise or counterclockwise.

## Swirling Diagrams: Alternative definitions



## Observation

A Diagram is swirling if and only if every arc is part of two swirls.

## Swirling Diagrams: Alternative definitions

## Observation

If in an arrangement of polygons all vertices are occluded, and each edge occludes vertices of at most one polygon, then the edges project into a swirling Diagram.


## Uniform Diagrams

Each arc in a Diagram feeds into exactly two arcs. So, the average number of arcs feeding into a given arc of a Diagram is two.


A Diagram is said uniform if each arc has two arcs feeding into it.

## Uniform Diagrams

## Proposition

All swirling Diagrams are uniform.


Proof. In a swirling Diagram, each arc is part of two distinct swirls, and so at least two arcs feed into it.

## Uniform Diagrams

## Proposition

All swirling Diagrams are uniform.


But each arc has two arcs feeding into it on average, so it must have exactly two arcs feeding into it.

## Uniform Diagrams



The converse is not true: there are uniform Diagrams that are not swirling.

## Uniform Diagrams



Note that the (portions of) arcs that are not part of a swirl form a cycle where each arc feeds into the next: this is not a coincidence...

## Uniform Diagrams

## Proposition

In a uniform Diagram, the non-swirling arcs form disjoint cycles.


Proof. Consider the last arc in a chain of non-swirling arcs.

## Uniform Diagrams

## Proposition

In a uniform Diagram, the non-swirling arcs form disjoint cycles.


This arc cannot form a swirl with the arc it feeds into (axiom 3).

## Uniform Diagrams

## Proposition

In a uniform Diagram, the non-swirling arcs form disjoint cycles.


So, the arc it feeds into cannot be part of two swirls (uniformity).

## Uniform Diagrams

## Proposition

In a uniform Diagram, the non-swirling arcs form disjoint cycles.


Therefore, the chain must be followed by another non-swirling arc.

## Uniform Diagrams

## Proposition

In a uniform Diagram, the non-swirling arcs form disjoint cycles.


Moreover, the chain can be uniquely extended backwards.

## Uniform Diagrams



Uniform Diagrams can have any number of unboundedly long cycles of non-swirling arcs.

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Uniform Diagrams can have any number of unboundedly long cycles of non-swirling arcs.

## Uniform Diagrams



By suitably merging consecutive arcs in each cycle, we can transform any uniform Diagram into a swirling one.

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## Swirl Graph



The Swirl Graph of a Diagram is an undirected multigraph on the set of swirls. For each arc shared by two swirls, there is an edge in the Swirl Graph.

## Swirl Graph

## Theorem

The Swirl Graph of a Diagram is a simple planar bipartite graph with non-empty partite sets.


Proof. Obviously the Swirl Graph is spherical, hence planar. The bipartition is given by the clockwise and counterclockwise swirls...

## Swirl Graph



Each edge in the Swirl Graph must connect a clockwise and a counterclockwise swirl. So the Swirl Graph is bipartite.

## Swirl Graph



To find a (counter)clockwise swirl, start anywhere and follow the Diagram (counter)clockwise. Hence the partite sets are not empty.

## Swirl Graph

Assume that the yellow swirl shares arcs $a$ and $b$ with another swirl. The second swirl must be located in the highlighted spherical lune.


Since $a$ goes upward, the second swirl must be in $A$. But $b$ goes downward, so the second swirl must be in $B$ : contradiction. Hence, the Swirl Graph is simple.

## Hemisphere lemma

## Lemma

In a Diagram, every hemisphere contains at least one full swirl.


Proof. Take any hemisphere $H$. Since the Diagram is connected and tiles are convex, there is an arc crossing the boundary of $H$.

## Hemisphere lemma

## Lemma

In a Diagram, every hemisphere contains at least one full swirl.


Follow the Diagram clockwise starting from this arc. If we remain in $H$, we eventually find a clockwise swirl fully contained in $H$.

## Hemisphere lemma

## Lemma

In a Diagram, every hemisphere contains at least one full swirl.


Otherwise, we find one arc whose clockwise endpoint is outside $H$. But then, the other endpoint is in $H$. Reach that endpoint.

## Hemisphere lemma

## Lemma

In a Diagram, every hemisphere contains at least one full swirl.


Continuing in this fashion, we either find a clockwise swirl in $H$, or eventually we enclose a region within $H$ by going counterclockwise.

## Hemisphere lemma

## Lemma

In a Diagram, every hemisphere contains at least one full swirl.


In this case, starting from the boundary of the enclosed region and following the Diagram counterclockwise, we eventually find a swirl.

## Hemisphere lemma



Note: In some cases, a hemisphere may contain exactly one swirl.

## More swirls

## Corollary

Every Diagram has at least 4 swirls.


Proof. We already know that a Diagram has 2 swirls.

## More swirls

## Corollary

Every Diagram has at least 4 swirls.


Take a great circle that properly intersects both swirls.

## More swirls

## Corollary

Every Diagram has at least 4 swirls.


By the previous lemma, each hemisphere contains one new swirl.

## Minimizing arcs (and swirls)

## Theorem

Every Diagram has at least 8 arcs, and there exist Diagrams with exactly 8 arcs (and exactly 4 swirls).


Proof. This is an example of a Diagram with 8 arcs and 4 swirls...

## Minimizing arcs (and swirls)



We know that a Diagram has at least 4 swirls. Obviously, if they do not share any arcs, then the Diagram has at least 12 arcs.

## Minimizing arcs (and swirls)



Since the Swirl Graph must be simple and bipartite, these 4 swirls can share at most 4 arcs. Thus the Diagram has at least 8 arcs.

## Minimizing visible edges

## Corollary

If a point does not see any vertices of a polyhedron, it sees at least 8 distinct edges. The bound is tight.


A lower-bound example can be constructed from this arrangement of 6 polygons, where the central point does not see any vertices and sees exactly 8 edges.

## Summary

- Spherical Occlusion Diagrams occur naturally when studying points that see no vertices of a polyhedron.
- There is a straightforward correspondence between swirling Diagrams and even-sided convex tilings of the sphere.
- Uniform Diagrams can be obtained by augmenting swirling Diagrams with disjoint non-swirling cycles.
- Swirls are patterns frequently appearing in Diagrams. By studying swirl graphs, we obtain a new Art Gallery theorem: If we see no vertices of a polyhedron, then we see 8+ edges.


## Future work

## Conjecture

There are no swirling Diagrams with $8,9,10,11,13,14,15,17$, 21, 22, 23, or 29 arcs.

## Conjecture

Every Diagram has at least 2 clockwise and 2 counterclockwise swirls.

## Conjecture

Not every Diagram is the projection of a polyhedron's 1-skeleton.

## Conjecture

Not every Diagram is combinatorially equivalent to the projection of a polyhedron's 1-skeleton.

## Conjecture

If a point does not see any vertices of a polyhedron, it sees at least 8 distinct faces.

## Spherical Occlusion Diagrams in everyday life



Modular origami: kusudama

## Spherical Occlusion Diagrams in everyday life



Modular origami: penultimate dodecahedron

## Spherical Occlusion Diagrams in everyday life



Modular origami: penultimate truncated icosahedron

## Spherical Occlusion Diagrams in everyday life

Kirigami ball decoration

## Spherical Occlusion Diagrams in everyday life



Monkey's fist knot

## Spherical Occlusion Diagrams in everyday life



Single-thread globe knot

## Spherical Occlusion Diagrams in everyday life



Double-thread globe knot

## Spherical Occlusion Diagrams in everyday life



Herringbone pineapple knot

## Spherical Occlusion Diagrams in everyday life



Stainless-steel globe knot

## Spherical Occlusion Diagrams in everyday life



Sepak-takraw ball

## Spherical Occlusion Diagrams in everyday life



Rattan balls

## Spherical Occlusion Diagrams in everyday life



Rattan vase

## Spherical Occlusion Diagrams in everyday life



Toroidal Occlusion Diagrams...?

