# A Theory of Spherical Diagrams CCCG 2022

Giovanni Viglietta

#### Japan Advanced Institute of Science and Technology (JAIST)

Toronto – August 27, 2022

# 3D Art Gallery Problem

Given a polyhedron in  $\mathbb{R}^3$ , choose a (preferably small) set of <u>vertices</u> or edges that collectively see its whole interior.



These are called **vertex guards** and **edge guards**.

# Vertex-guarding polyhedra

The Art Gallery Problem for *vertex guards* may be <u>unsolvable</u>, even in some orthogonal polyhedra:



Some points in the central region are invisible to all vertices!

# Point-guarding polyhedra

Even if guards are not constrained to lie on vertices, there are (orthogonal) polyhedra that require  $\Omega(n\sqrt{n})$  point guards!







# Edge-guarding polyhedra

These observations justify the study of *edge guards*.



Problem 1. How many edge guards are needed for a polyhedron?Problem 2. Assuming that there is a point guard on every vertex of a polyhedron, how many additional edge guards are needed?



When polygons in  $\mathbb{R}^3$  are orthographically projected onto a sphere, their edges become arcs of great circles.



Moreover, when a polygon is <u>partially hidden</u> (i.e., **"occluded"**) by another, in the projection there are arcs feeding into other arcs.



Moreover, when a polygon is <u>partially hidden</u> (i.e., **"occluded"**) by another, in the projection there are arcs feeding into other arcs.



If in an arrangement of polygons <u>all vertices are occluded</u>, then their edges project into a **"Spherical Occlusion Diagram"**.



If in an arrangement of polygons <u>all vertices are occluded</u>, then their edges project into a **"Spherical Occlusion Diagram"**.



If in an arrangement of polygons <u>all vertices are occluded</u>, then their edges project into a **"Spherical Occlusion Diagram"**.



In particular, this applies to polyhedra: if all vertices are occluded, then the 1-skeleton projects into a Spherical Occlusion Diagram.



A **Spherical Occlusion Diagram**, or just "Diagram", is a finite non-empty collection of arcs of great circles on the unit sphere.



All arcs in a Diagram must be internally disjoint.



Both endpoints of each arc in a Diagram must lie in the interiors of some other arcs in the Diagram (every arc "feeds into" two arcs).



All the arcs in a Diagram that feed into the same arc must reach it from the <u>same side</u>.



All the arcs in a Diagram that feed into the same arc must reach it from the <u>same side</u>.

# Spherical Occlusion Diagrams: Examples

#### **Diagram axioms:**

- 1. Arcs are internally disjoint.
- 2. Each arc feeds into two arcs.
- 3. All arcs that feed into the same arc reach it from the same side.

# Spherical Occlusion Diagrams: Examples

### Diagram axioms:

- 1. Arcs are internally disjoint.
- 2. Each arc feeds into two arcs.
- 3. All arcs that feed into the same arc reach it from the same side.

# Spherical Occlusion Diagrams: Examples

### Diagram axioms:

- 1. Arcs are internally disjoint.
- 2. Each arc feeds into two arcs.
- 3. All arcs that feed into the same arc reach it from the same side.

#### Proposition

No arc in a Diagram is longer than a great semicircle.



Proof. Otherwise it would have arcs feeding into it from both sides.

#### Proposition

No arc in a Diagram is longer than a great semicircle.



Proof. Otherwise it would have arcs feeding into it from both sides.

#### Proposition

No arc in a Diagram is longer than a great semicircle.



Proof. Otherwise it would have arcs feeding into it from both sides.

#### Corollary

No two arcs in a Diagram feed into each other.



**Proof.** Otherwise they would be longer than a great semicircle.

### Proposition

A Diagram partitions the sphere into convex regions (or "tiles").



**Proof.** Two points in the same region can be connected by a chain of arcs of great circles that does not intersect the Diagram.

### Proposition

A Diagram partitions the sphere into convex regions (or "tiles").



The arc joining the first and the third vertex of the chain do not intersect the Diagram, either...

### Proposition

A Diagram partitions the sphere into convex regions (or "tiles").



...Otherwise, following the Diagram we would intersect the first two arcs in the chain, which is impossible by assumption.

### Proposition

A Diagram partitions the sphere into convex regions (or "tiles").



So we can simplify the chain, reducing it by one arc. Inductively repeating this reasoning, we can reduce the chain to a single arc.

### Proposition

A Diagram partitions the sphere into convex regions (or "tiles").



Any two points in the region are connected by an arc of a great circle that does not intersect the Diagram; hence, it is convex.

#### Corollary

Every Diagram is connected.



**Proof.** If there are two connected components, each of them is a Diagram. So, one is contained in a tile  $\mathcal{F}$  determined by the other.

#### Corollary

Every Diagram is connected.



Take an arc in  $\mathcal{F}$  with endpoints close to the first component that intersects the second component.

#### Corollary

Every Diagram is connected.



The arc can be replaced by a chain that intersects neither connected component of the Diagram.

#### Corollary

Every Diagram is connected.



So its endpoints are in the same tile determined by the whole Diagram, and this tile cannot be convex.

#### Proposition

A Diagram with n arcs partitions the sphere into n + 2 tiles.



**Proof.** A Diagram induces a planar graph with v vertices and n + v edges. By Euler's formula, f + v = n + v + 2, hence f = n + 2.



How can we automatically generate large classes of Diagrams?



Start from a subdivision of the sphere into strictly <u>convex tiles</u>, where each tile has an even number of edges.
#### Spherical Occlusion Diagrams: Construction



Note that the 1-skeleton of the tiling is <u>bipartite</u>, because it has no odd cycles.

#### Spherical Occlusion Diagrams: Construction



We can turn each vertex of the tiling into a "**swirl**" going clockwise or counterclockwise according to the bipartition of the 1-skeleton.

#### Spherical Occlusion Diagrams: Construction



This operation defines a natural correspondence between even-sided spherical tilings and so-called *swirling Diagrams*.

This method enables the <u>automatic construction</u> of swirling Diagrams from convex tilings of the sphere or convex polyhedra.



Prisms with even-sided bases

This method enables the <u>automatic construction</u> of swirling Diagrams from convex tilings of the sphere or convex polyhedra.



Truncated antiprisms

This method enables the <u>automatic construction</u> of swirling Diagrams from convex tilings of the sphere or convex polyhedra.



Truncated bipyramids with even-degree vertices

This method enables the <u>automatic construction</u> of swirling Diagrams from convex tilings of the sphere or convex polyhedra.



Trapezohedra

This method enables the <u>automatic construction</u> of swirling Diagrams from convex tilings of the sphere or convex polyhedra.



Rhombic dodecahedron

This method enables the <u>automatic construction</u> of swirling Diagrams from convex tilings of the sphere or convex polyhedra.



Deltoidal icositetrahedron

This method enables the <u>automatic construction</u> of swirling Diagrams from convex tilings of the sphere or convex polyhedra.



Rhombic triancontahedron

This method enables the <u>automatic construction</u> of swirling Diagrams from convex tilings of the sphere or convex polyhedra.



Deltoidal hexecontahedron

This method enables the <u>automatic construction</u> of swirling Diagrams from convex tilings of the sphere or convex polyhedra.



Truncated cuboctahedron

This method enables the <u>automatic construction</u> of swirling Diagrams from convex tilings of the sphere or convex polyhedra.



Truncated icosidodecahedron

#### Swirling Diagrams: Alternative definitions



As we saw, a **swirl** in a Diagram is a cycle of arcs such that each arc feeds into the next going clockwise or counterclockwise.

#### Swirling Diagrams: Alternative definitions



#### Observation

A Diagram is swirling if and only if every arc is part of two swirls.

#### Swirling Diagrams: Alternative definitions

#### Observation

If in an arrangement of polygons all vertices are occluded, and each edge occludes vertices of at most one polygon, then the edges project into a swirling Diagram.



Each arc in a Diagram feeds into exactly <u>two arcs</u>. So, the average number of arcs feeding into a given arc of a Diagram is two.



A Diagram is said uniform if each arc has two arcs feeding into it.

#### Proposition

All swirling Diagrams are uniform.



**Proof.** In a swirling Diagram, each arc is part of two distinct swirls, and so <u>at least two arcs</u> feed into it.

#### Proposition

All swirling Diagrams are uniform.



But each arc has two arcs feeding into it <u>on average</u>, so it must have exactly two arcs feeding into it.



The converse is not true: there are  $\underline{\text{uniform Diagrams}}$  that are not swirling.



Note that the (portions of) arcs that are not part of a swirl form a cycle where each arc feeds into the next: *this is not a coincidence...* 

#### Proposition

In a uniform Diagram, the non-swirling arcs form disjoint cycles.



**Proof.** Consider the last arc in a chain of non-swirling arcs.

#### Proposition

In a uniform Diagram, the non-swirling arcs form disjoint cycles.



This arc cannot form a swirl with the arc it feeds into (axiom 3).

#### Proposition

In a uniform Diagram, the non-swirling arcs form disjoint cycles.



So, the arc it feeds into cannot be part of two swirls (uniformity).

#### Proposition

In a uniform Diagram, the non-swirling arcs form disjoint cycles.



Therefore, the chain must be followed by another non-swirling arc.

#### Proposition

In a uniform Diagram, the non-swirling arcs form disjoint cycles.



Moreover, the chain can be uniquely extended backwards.



Uniform Diagrams can have <u>any number</u> of <u>unboundedly long</u> cycles of non-swirling arcs.



Uniform Diagrams can have <u>any number</u> of <u>unboundedly long</u> cycles of non-swirling arcs.













The **Swirl Graph** of a Diagram is an <u>undirected multigraph</u> on the set of swirls. For each arc shared by two swirls, there is an edge in the Swirl Graph.

# Swirl Graph

#### Theorem

The Swirl Graph of a Diagram is a simple planar bipartite graph with non-empty partite sets.



**Proof.** Obviously the Swirl Graph is spherical, hence planar. The bipartition is given by the clockwise and counterclockwise swirls...



Each edge in the Swirl Graph must connect a clockwise and a counterclockwise swirl. So the Swirl Graph is bipartite.


To find a (counter)clockwise swirl, start anywhere and follow the Diagram (counter)clockwise. Hence the partite sets are not empty.

# Swirl Graph

Assume that the yellow swirl shares arcs a and b with another swirl. The second swirl must be located in the highlighted spherical lune.



Since a goes upward, the second swirl must be in A. But b goes downward, so the second swirl must be in B: contradiction. Hence, the Swirl Graph is simple.

#### Lemma

In a Diagram, every hemisphere contains at least one full swirl.



**Proof.** Take any hemisphere H. Since the Diagram is connected and tiles are convex, there is an arc crossing the boundary of H.

#### Lemma

In a Diagram, every hemisphere contains at least one full swirl.



Follow the Diagram clockwise starting from this arc. If we remain in H, we eventually find a clockwise swirl fully contained in H.

#### Lemma

In a Diagram, every hemisphere contains at least one full swirl.



Otherwise, we find one arc whose clockwise endpoint is outside H. But then, the other endpoint is in H. Reach that endpoint.

#### Lemma

In a Diagram, every hemisphere contains at least one full swirl.



Continuing in this fashion, we either find a clockwise swirl in H, or eventually we enclose a region within H by going counterclockwise.

#### Lemma

In a Diagram, every hemisphere contains at least one full swirl.



In this case, starting from the boundary of the enclosed region and following the Diagram counterclockwise, we eventually find a swirl.



Note: In some cases, a hemisphere may contain exactly one swirl.

## More swirls

## Corollary

Every Diagram has at least 4 swirls.



**Proof.** We already know that a Diagram has 2 swirls.

# More swirls

## Corollary

Every Diagram has at least 4 swirls.



Take a great circle that properly intersects both swirls.

# More swirls

## Corollary

Every Diagram has at least 4 swirls.



By the previous lemma, each hemisphere contains one new swirl.

# Minimizing arcs (and swirls)

#### Theorem

Every Diagram has <u>at least 8 arcs</u>, and there exist Diagrams with exactly 8 arcs (and exactly 4 swirls).



Proof. This is an example of a Diagram with 8 arcs and 4 swirls...

# Minimizing arcs (and swirls)



We know that a Diagram has at least 4 swirls. Obviously, if they do not share any arcs, then the Diagram has at least 12 arcs.

# Minimizing arcs (and swirls)



Since the Swirl Graph must be simple and bipartite, these 4 swirls can share at most 4 arcs. Thus the Diagram has at least 8 arcs.

# Minimizing visible edges

### Corollary

If a point does not see any vertices of a polyhedron, it sees at least 8 distinct edges. The bound is tight.



A lower-bound example can be constructed from this arrangement of 6 polygons, where the central point does not see any vertices and sees exactly 8 edges.

- **Spherical Occlusion Diagrams** occur naturally when studying points that see no vertices of a polyhedron.
- There is a straightforward correspondence between **swirling Diagrams** and even-sided convex tilings of the sphere.
- **Uniform Diagrams** can be obtained by augmenting swirling Diagrams with disjoint non-swirling cycles.
- Swirls are patterns frequently appearing in Diagrams. By studying **swirl graphs**, we obtain a new Art Gallery theorem: *If we see no vertices of a polyhedron, then we see 8+ edges.*

# Future work

### Conjecture

There are no swirling Diagrams with 8, 9, 10, 11, 13, 14, 15, 17, 21, 22, 23, or 29 arcs.

#### Conjecture

Every Diagram has at least 2 clockwise and 2 counterclockwise swirls.

## Conjecture

Not every Diagram is the projection of a polyhedron's 1-skeleton.

#### Conjecture

Not every Diagram is <u>combinatorially equivalent</u> to the projection of a polyhedron's 1-skeleton.

### Conjecture

If a point does not see any vertices of a polyhedron, it sees at least 8 distinct <u>faces</u>.



Modular origami: kusudama



Modular origami: penultimate dodecahedron



Modular origami: penultimate truncated icosahedron



Kirigami ball decoration



Monkey's fist knot



Single-thread globe knot



Double-thread globe knot



Herringbone pineapple knot



Stainless-steel globe knot



Sepak-takraw ball



Rattan balls



Rattan vase



Toroidal Occlusion Diagrams...?