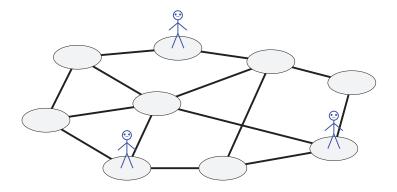
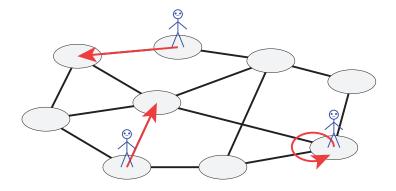
# Universal Systems of Oblivious Mobile Robots SIROCCO 2016

Paola Flocchini, Nicola Santoro, Giovanni Viglietta, Masafumi Yamashita

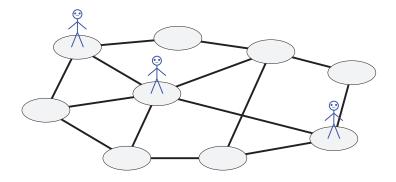
Helsinki - July 20, 2016



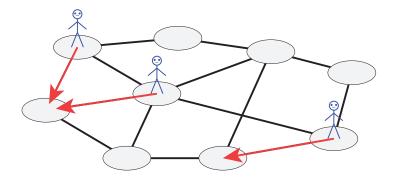
We consider a set of anonymous robots living on a network.



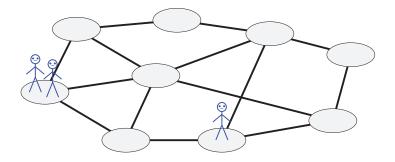
At each *round*, all robots move simultaneously to an adjacent vertex (or stay still).



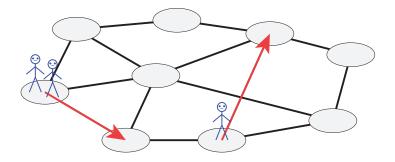
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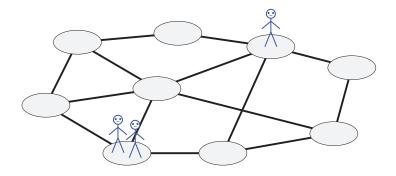
A robot's destination is computed via a deterministic algorithm whose input is only the current configuration (*obliviousness*).



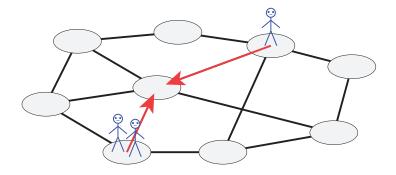
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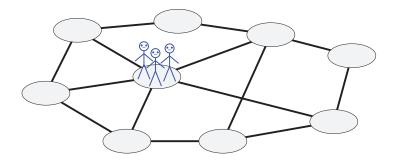
Several robots may occupy the same vertex.



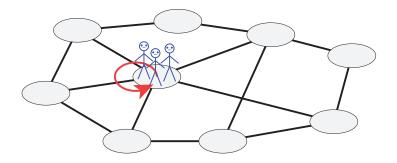
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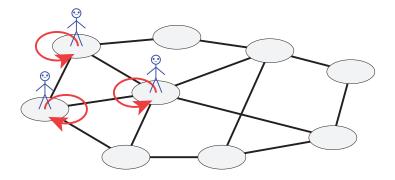
A basic well-studied problem is gathering.



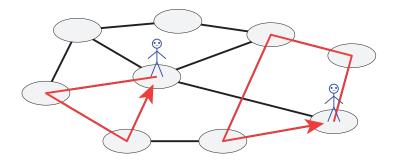
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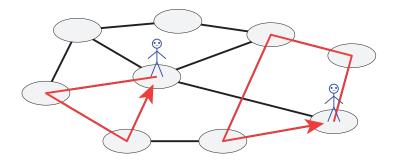
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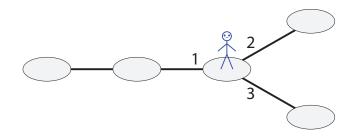
A more general problem is *pattern formation* (e.g., a triangle).



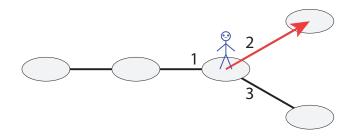
Another common task may be to implement a self-stabilizing *clock* with a specific period (e.g., 12).



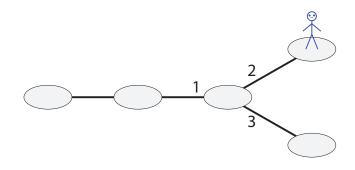
Our goal is to define what it means for such a system to *compute a function*, and determine what functions it can compute.



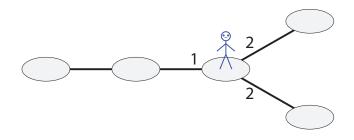
#### Each vertex of the network has port labels on incident edges.



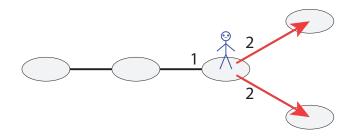
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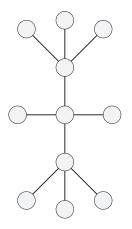
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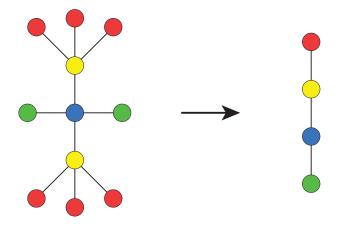
Labels need not be unique. This may cause non-determinism in executions (even though algorithms are deterministic!).



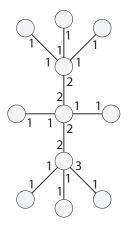
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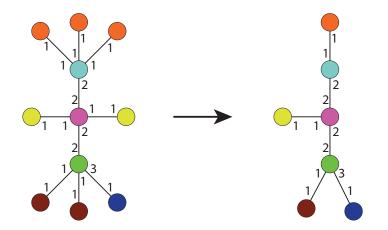
Identifying equivalent vertices yields the quotient graph.



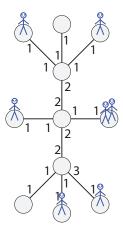
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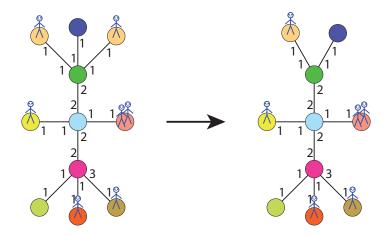
Vertices can also be distinguished by their surrounding labels...



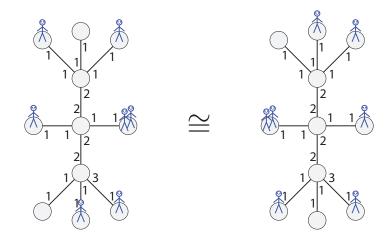
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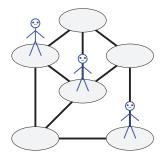
...And by the amount of robots that occupy them.



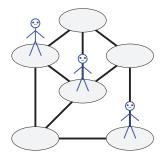
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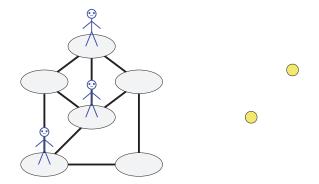
Vertices and robots are otherwise *indistinguishable*.



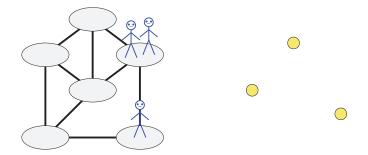
#### Non-equivalent configurations can be arranged in a graph.



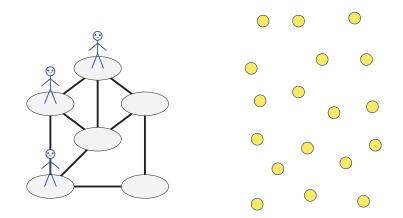
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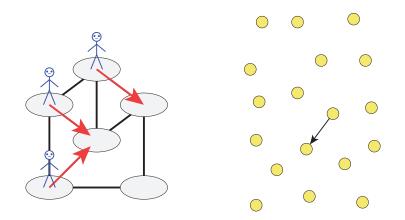
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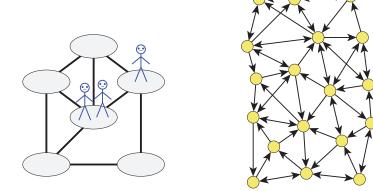
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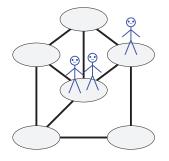
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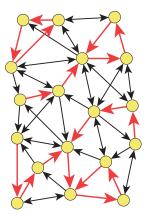


An edge in the graph means that two configurations are connected by a simultaneous move of all the robots.

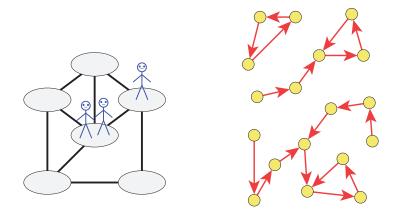


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An *algorithm* is a selection of an outgoing edge for each vertex of the *configuration graph*.

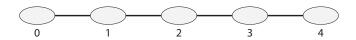


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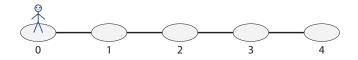
## Example: oriented path

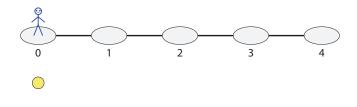


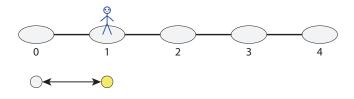
Let us study the configuration graph of a *path* with left-right *orientation*.

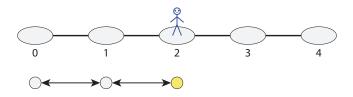


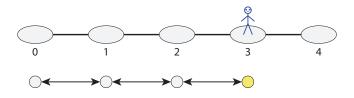
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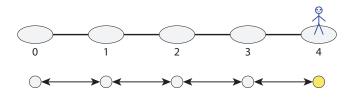


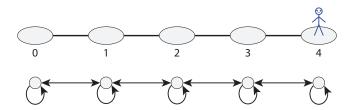


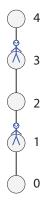




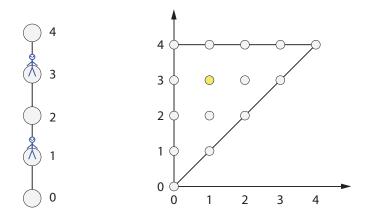




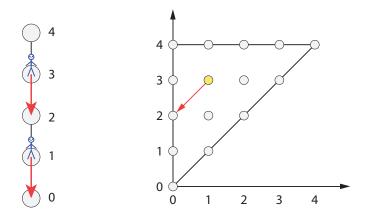




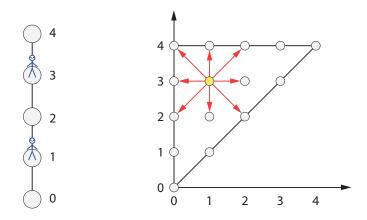
If there are two robots, the configuration graph is a triangular mesh.



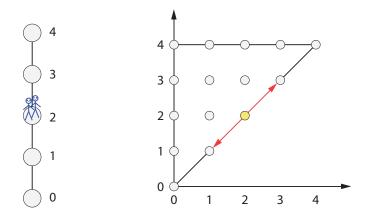
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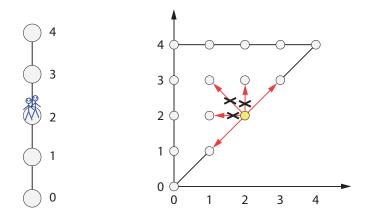
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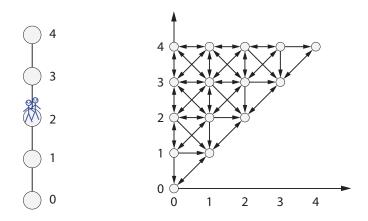
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When the two robots meet, they can no longer be separated.



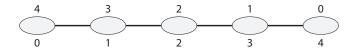
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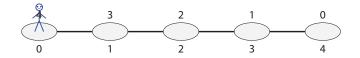
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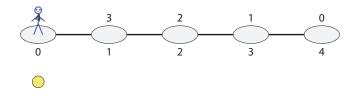


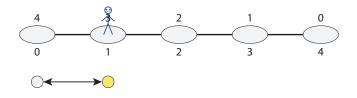
Suppose now the path is unoriented.

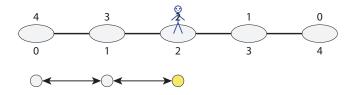


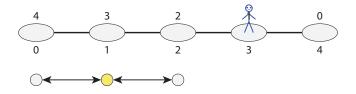
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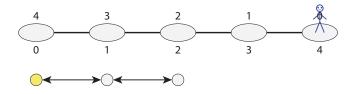


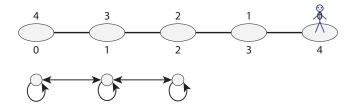


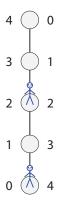




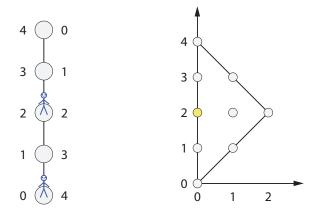




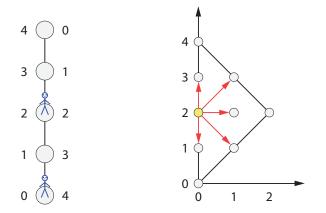




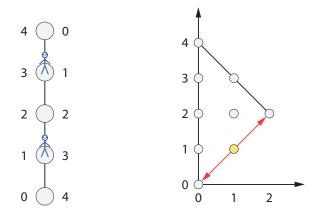
If there are two robots, the configuration graph is a triangular mesh of roughly half the size.



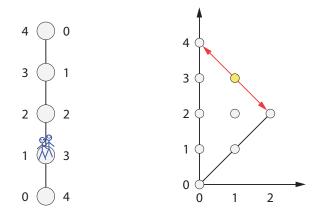
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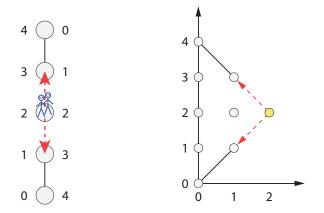
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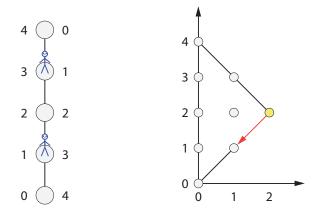


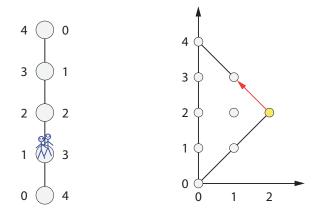
If the robots are in distinct symmetric locations, they must move symmetrically.

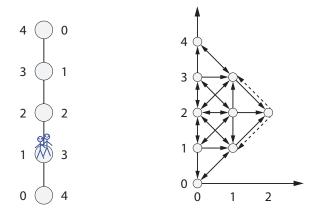


If the robots are on a same non-central vertex, they must remain together.

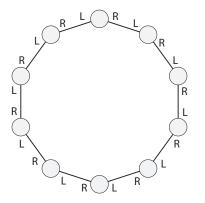






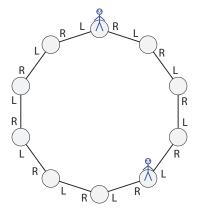


# Example: oriented ring



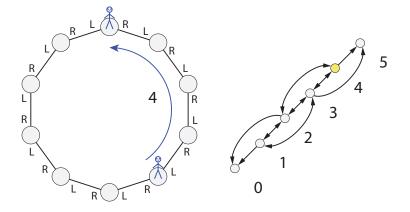
Let the network be a *ring* with left-right orientation.

### Example: oriented ring



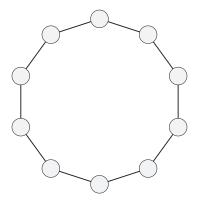
If two robots are on it, the configuration is fully determined by their distance.

### Example: oriented ring



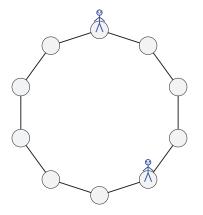
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#### Example: unoriented ring



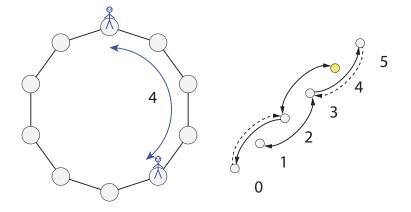
Let the network be an unoriented ring.

#### Example: unoriented ring

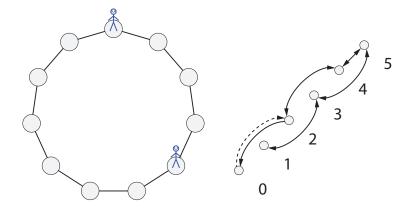


A configuration of two robots is again determined by their distance, but now fewer moves are possible.

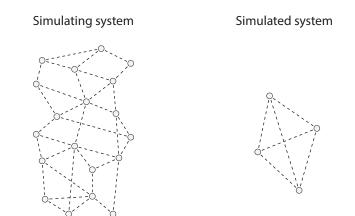
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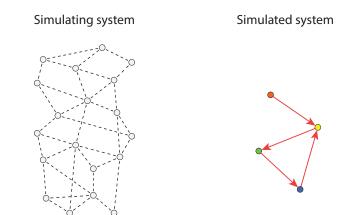
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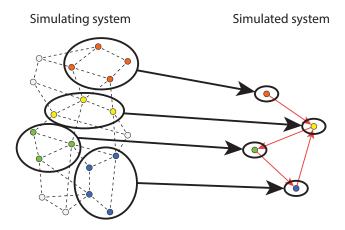
If the network's size is even, the configuration graph consists of two independent paths; if the size is odd, it consists of a single path.



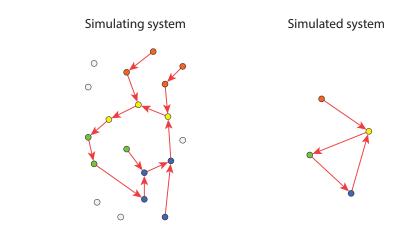
What does it mean for an *algorithm* for a given system (network, swarm of robots) to *simulate* an algorithm for another system?



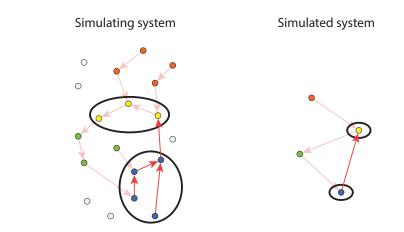
What does it mean for an *algorithm* for a given system (network, swarm of robots) to *simulate* an algorithm for another system?



Each configuration of the simulated system must be *represented* by a set of configurations of the simulating system...



...In such a way that the simulating algorithm "respects" this correspondence.



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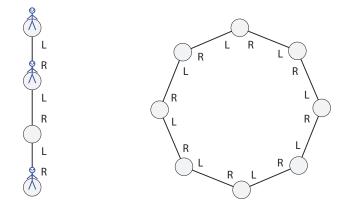


Any algorithm for an oriented n-path can be simulated on an unoriented 2n-path by the same number of robots.

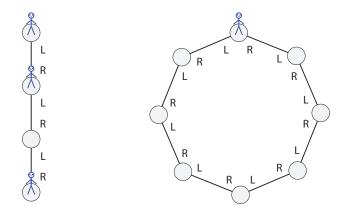




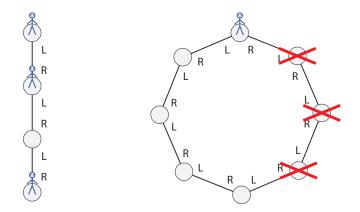
The simulation takes place on one half of the path, while the other half remains empty and defines an implicit orientation.



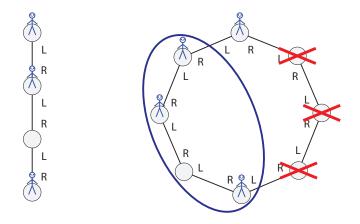
Any algorithm for an oriented n-path can be simulated on an oriented ring of size 2n by adding one robot.



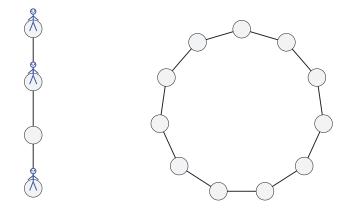
One half of the ring consists of a stationary *pivot robot* followed by n - 1 empty vertices.



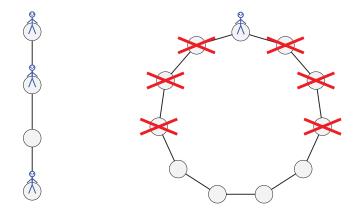
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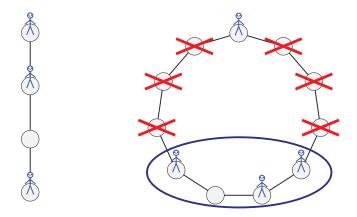
The simulation takes place on the remaining vertices, which are uniquely identified.



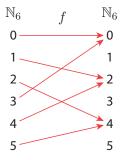
Any algorithm for an unoriented *n*-path can be simulated on a larger unoriented ring of size 3n - 1 by adding one robot.



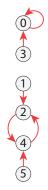
Place a static *pivot robot* surrounded on both sides by n - 1 empty vertices.



The simulation takes place on the remaining n vertices, which are uniquely determined.



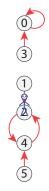
What does it mean to *compute* a function  $f : \mathbb{N}_k \to \mathbb{N}_k$ ?



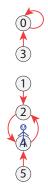
Interpret the function as a graph on  $\mathbb{N}_{k}$ ...



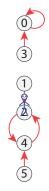
...Representing an algorithm for one robot.



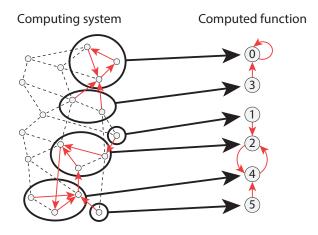
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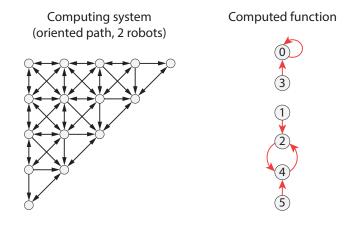
...Representing an algorithm for one robot.



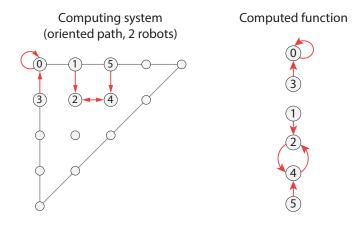
...Representing an algorithm for one robot.



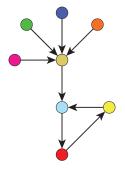
Computing the function means to simulate that algorithm.



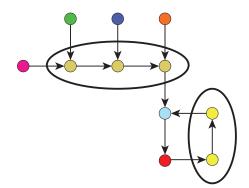
The computing system could be a long-enough oriented path with 2 robots.



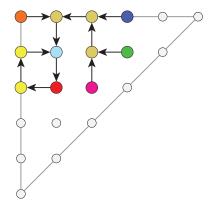
Indeed, the function's graph can be drawn on a grid...



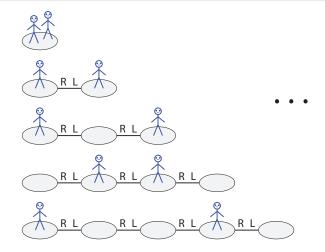
...Provided that we reduce the maximum degree to 3 and we make the length of every cycle even, by introducing extra vertices.



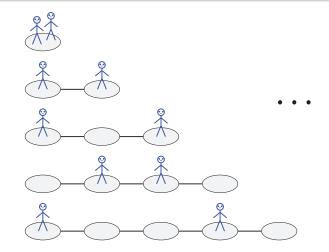
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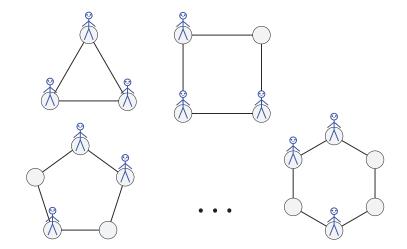
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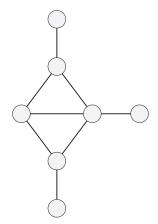
Therefore, the set of all oriented paths with 2 robots is *universal*, in that it can compute all (finite) functions.



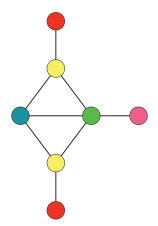
Since unoriented paths can simulate oriented paths, also the set of all unoriented paths with 2 robots is universal.



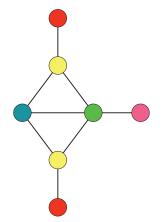
Similarly, the set of all (oriented or unoriented) rings with 3 robots is universal.

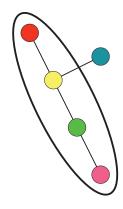


Suppose that a network's quotient graph contains a sub-path of length, say, 4.

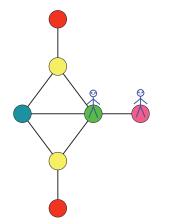


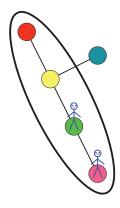
Suppose that a network's quotient graph contains a sub-path of length, say, 4.

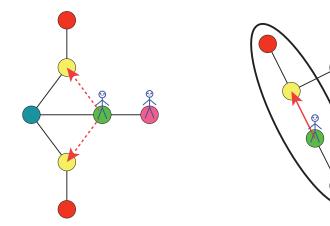


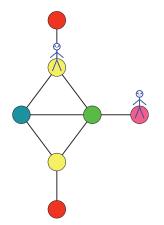


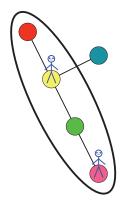
Suppose that a network's quotient graph contains a sub-path of length, say, 4.

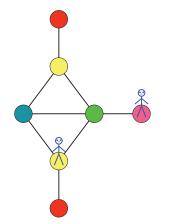


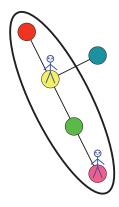


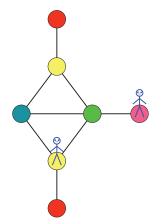


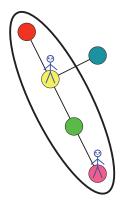




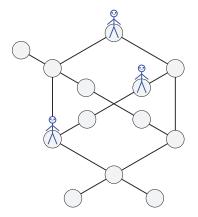






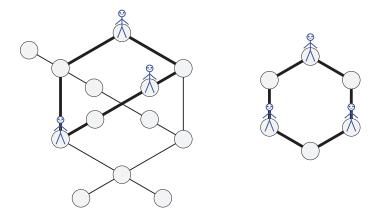


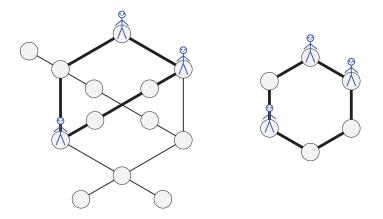
So, any set of networks with 2 robots whose quotient graphs contain arbitrarily long sub-paths is universal.

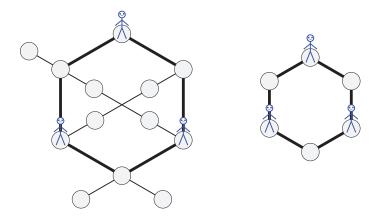


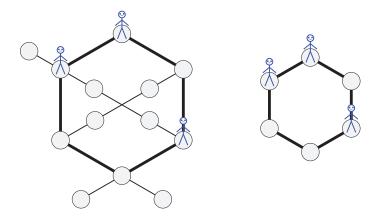
Suppose now that a network's girth is, say, 6.

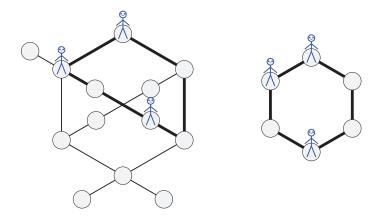
Universal Systems of Oblivious Mobile Robots

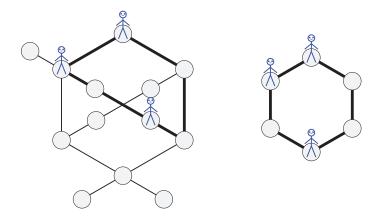












So, any set of networks of arbitrarily large (finite) girths and 3 robots is universal.

#### Theorem

Any set of networks with 2 robots whose quotient graphs contain unboundedly long sub-paths is universal.

#### Theorem

Any set of networks with 3 robots and unboundedly large (finite) girths is universal.

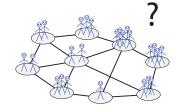
#### Conjecture

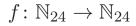
A set of anonymous networks with 3 robots is universal if and only if they have unboundedly large (finite) girths or their quotient graphs contain unboundedly long sub-paths.

$$f: \mathbb{N}_{24} \to \mathbb{N}_{24}$$

Suppose we wanted to compute all the functions on a given set, say,  $\mathbb{N}_{24}.$ 

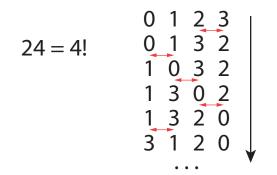
Universal Systems of Oblivious Mobile Robots



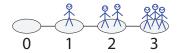


What is the *smallest* network on which we can compute all of them (possibly using a large number of robots)?

Universal Systems of Oblivious Mobile Robots



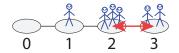
Recall that the 24 permutations of 4 objects can be ordered in such a way that two consecutive permutations differ by a transposition of two adjacent objects.



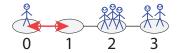


So, by putting a different number of robots on each vertex of a 4-path, we can encode the position of a robot on a 24-path...

Universal Systems of Oblivious Mobile Robots



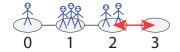




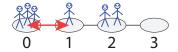




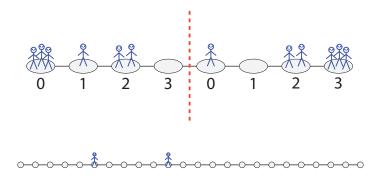




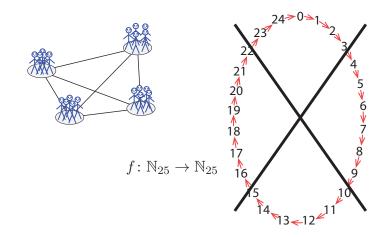




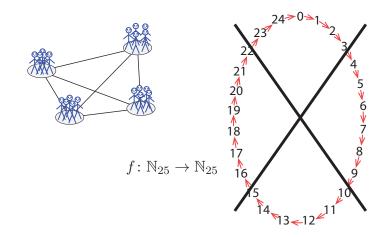




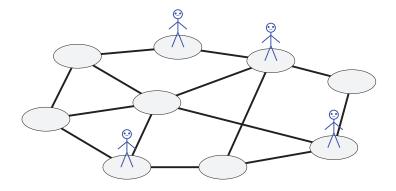
With an 8-path, we can encode the positions of 2 robots on a 24-path, and hence compute any function  $f: \mathbb{N}_{24} \to \mathbb{N}_{24}$ .

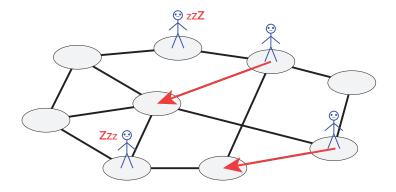


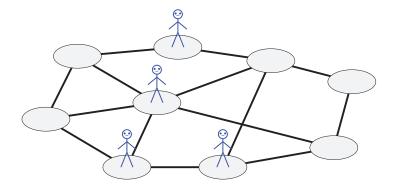
On the other hand, on a network of size 4 we cannot compute a function whose graph is a cycle of length 25.

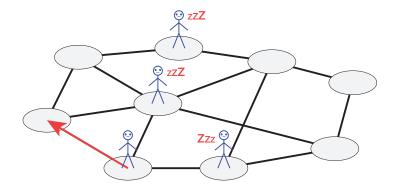


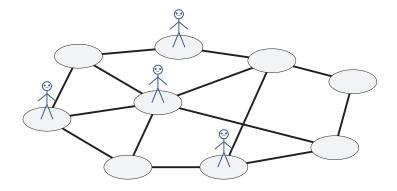
Therefore, we can determine the optimal size up to a factor of 2 (such is the ratio of our upper and lower bounds).

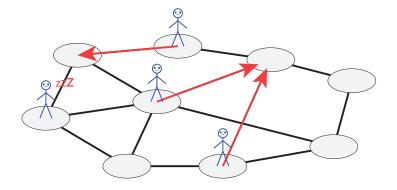












We can compare the "power" of the two models by comparing the functions that can be computed in each of them.