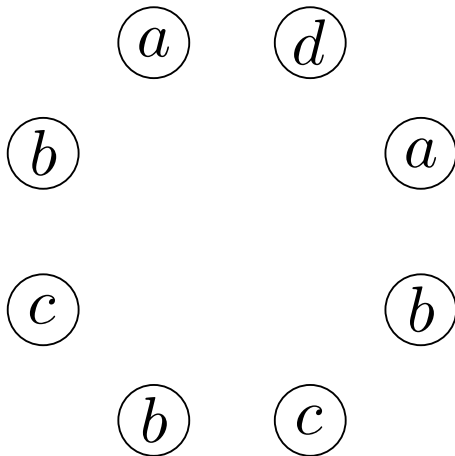


On the Power of Weaker Pairwise Interaction:
Fault-Tolerant Simulation of Population Protocols
ICDCS 2017

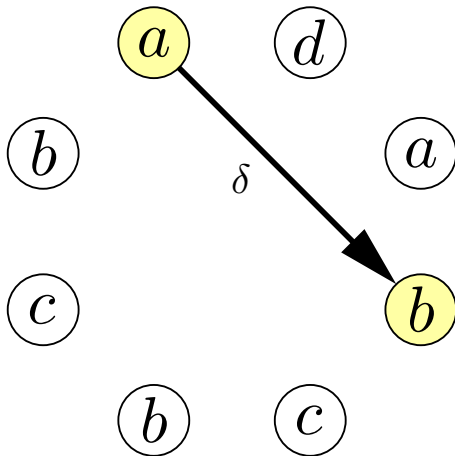
Giuseppe Antonio Di Luna, Paola Flocchini, Taisuke Izumi,
Tomoko Izumi, Nicola Santoro, Giovanni Viglietta

Atlanta – June 8, 2017



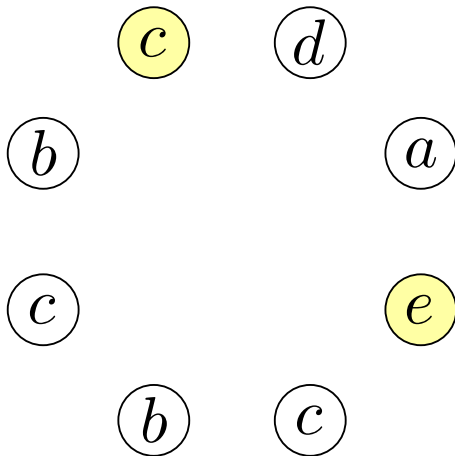
Setting: a set of finite-state agents.

Population Protocols

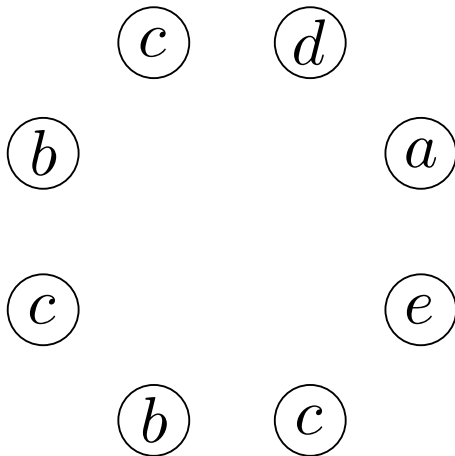


Pairs of agents interact in a non-deterministic order...

Population Protocols

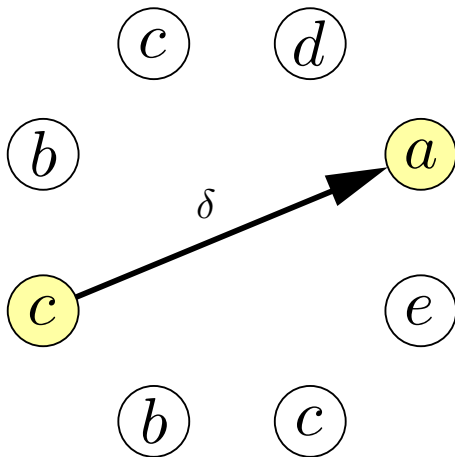


...and change states according to a transition function.



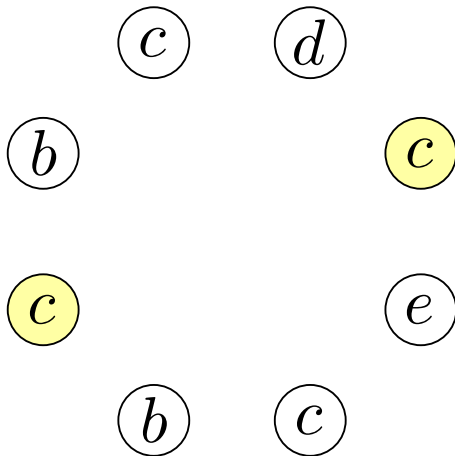
...and change states according to a transition function.

Population Protocols

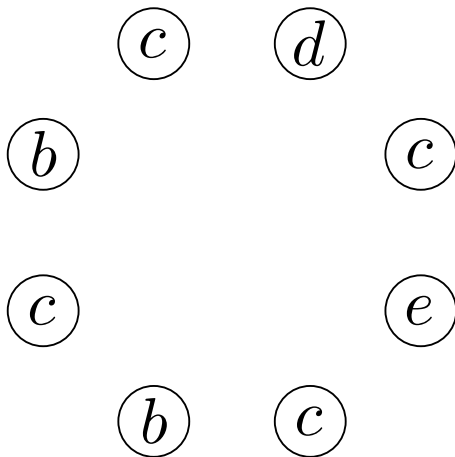


...and change states according to a transition function.

Population Protocols

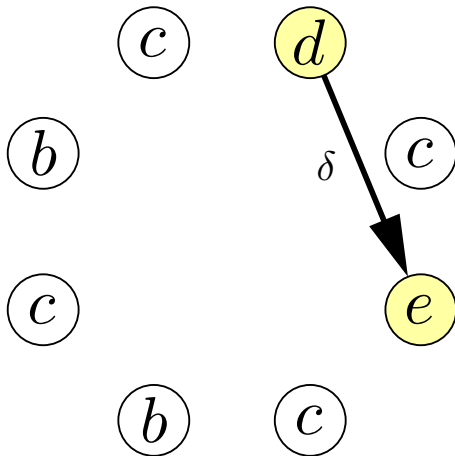


...and change states according to a transition function.



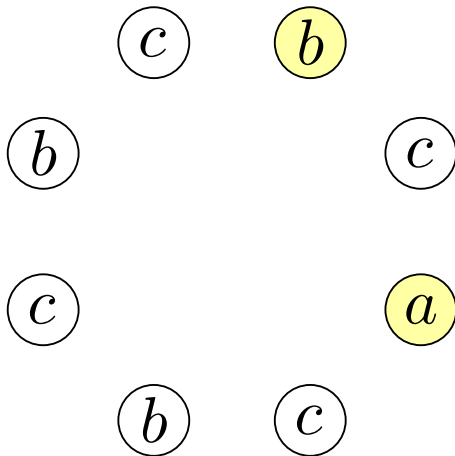
...and change states according to a transition function.

Population Protocols



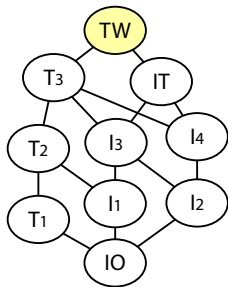
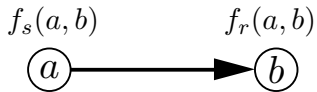
...and change states according to a transition function.

Population Protocols



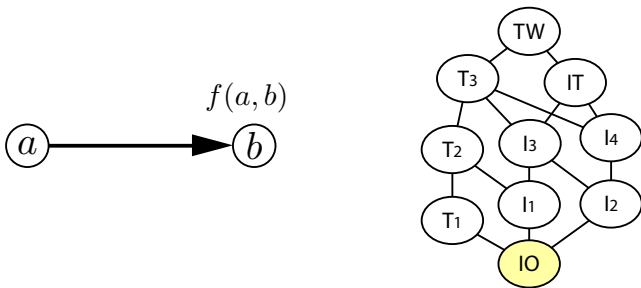
...and change states according to a transition function.

One-way models and omission faults



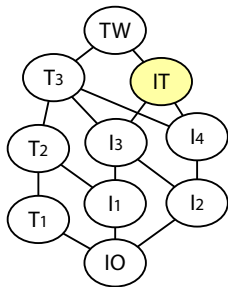
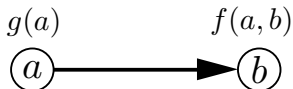
The traditional interaction model is called **Two-Way**.

One-way models and omission faults



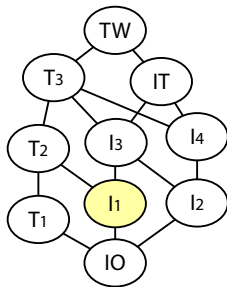
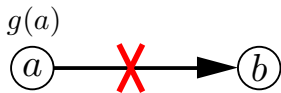
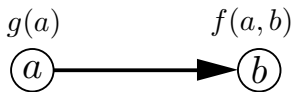
Immediate Observation: only the second agent transitions.

One-way models and omission faults



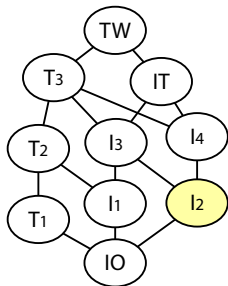
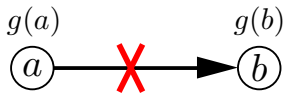
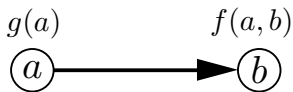
Immediate Transmission: the first agent detects *proximity*.

One-way models and omission faults



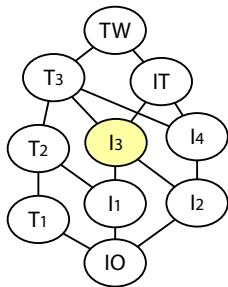
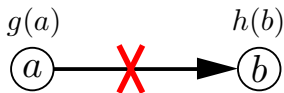
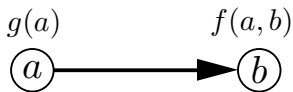
I_1 : IT with omission faults, no detection.

One-way models and omission faults



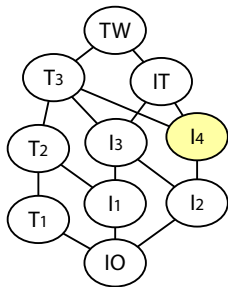
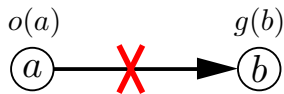
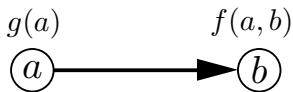
I₂: IT with omission faults, proximity detection.

One-way models and omission faults



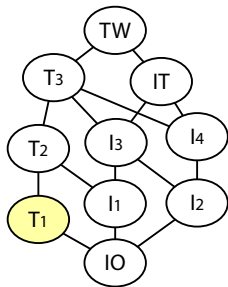
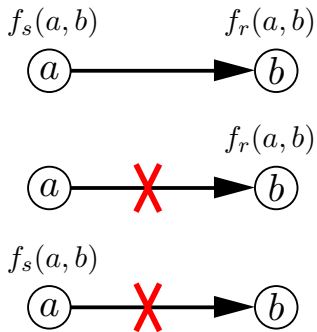
I₃: IT with omission faults, reactor-side omission detection.

One-way models and omission faults



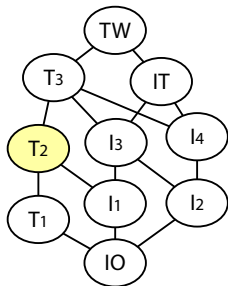
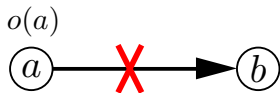
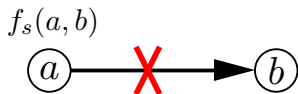
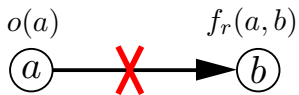
I₄: IT with omission faults, starter-side omission detection.

One-way models and omission faults



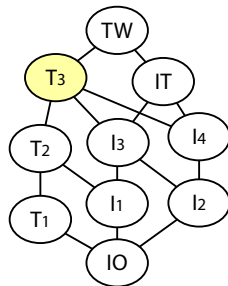
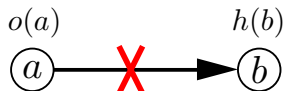
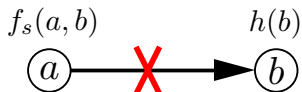
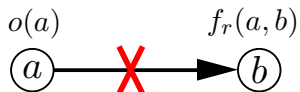
T_1 : TW with omission faults, no detection.

One-way models and omission faults



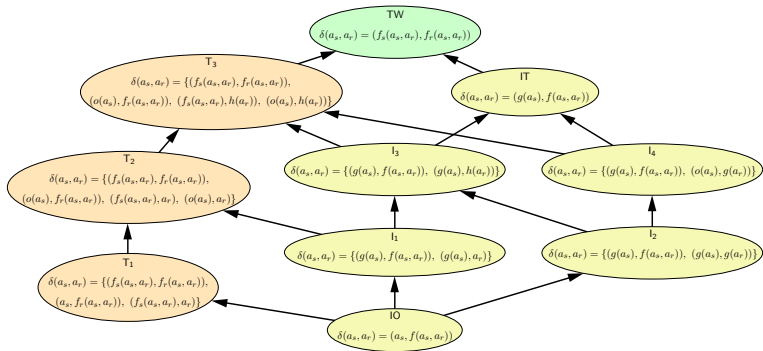
T₂: TW with omission faults, starter-side omission detection.

One-way models and omission faults



T₃: TW with omission faults, omission detection by both sides.

One-way models and omission faults




Theorem: all possible models obtained by combining one-way and two-way interactions with omission detection and proximity detection, starter-side or reactor-side, fall into one of these classes.

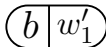
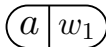
Simulating TW protocols with weaker ones



We seek to *simulate* two-way interactions in weaker models.

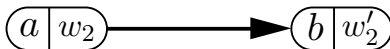
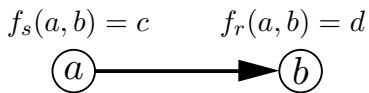
Simulating TW protocols with weaker ones

$$f_s(a, b) = c \quad f_r(a, b) = d$$





The simulating agents have a *simulated state* and a *work state*.

Simulating TW protocols with weaker ones



Typically, an interaction determines a change in the work state.


Simulating TW protocols with weaker ones

$$f_s(a, b) = c \quad f_r(a, b) = d$$




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Simulating TW protocols with weaker ones


$$f_s(a, b) = c \quad f_r(a, b) = d$$


A diagram showing a transition from state a to state b . The state a is represented by a circle containing the letter a , and the state b is represented by a circle containing the letter b . A thick black arrow points from the circle a to the circle b .



Typically, an interaction determines a change in the work state.

Simulating TW protocols with weaker ones


$$f_s(a, b) = c \quad f_r(a, b) = d$$


A diagram showing a transition from state a to state b . The state a is represented by a circle containing the letter a , and the state b is represented by a circle containing the letter b . A thick black arrow points from the circle a to the circle b .



Occasionally, changes in the simulated state may occur.


Simulating TW protocols with weaker ones

$$f_s(a, b) = c \quad f_r(a, b) = d$$


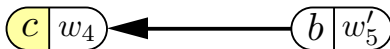


Occasionally, changes in the simulated state may occur.

Simulating TW protocols with weaker ones


$$f_s(a, b) = c \quad f_r(a, b) = d$$


A diagram showing a transition from state a to state b . State a is represented by a circle containing the letter a . A thick black arrow points from state a to state b , which is represented by a circle containing the letter b .

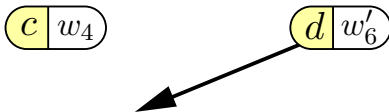


Occasionally, changes in the simulated state may occur.

Simulating TW protocols with weaker ones

$$f_s(a, b) = c \quad f_r(a, b) = d$$


A diagram showing a transition from node a to node b . Node a is a circle containing the letter a , and node b is a circle containing the letter b . A thick black arrow points from node a to node b .



These have to mimic transitions in the simulated TW protocol.

Simulating TW protocols with weaker ones

c

d

$c \mid w_4$

$d \mid w'_6$

Globally, we want to pair up simulated states transitions...

Simulating TW protocols with weaker ones

c

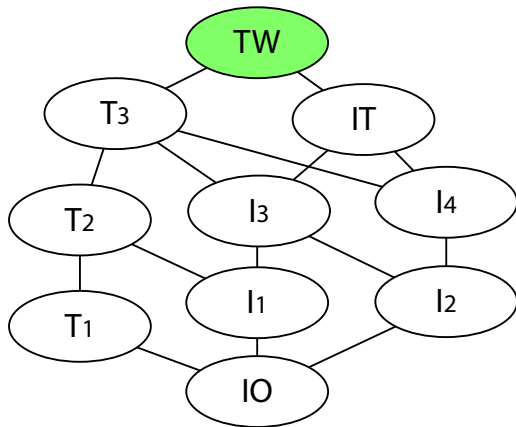
d

$c \mid w_4$

$d \mid w'_6$

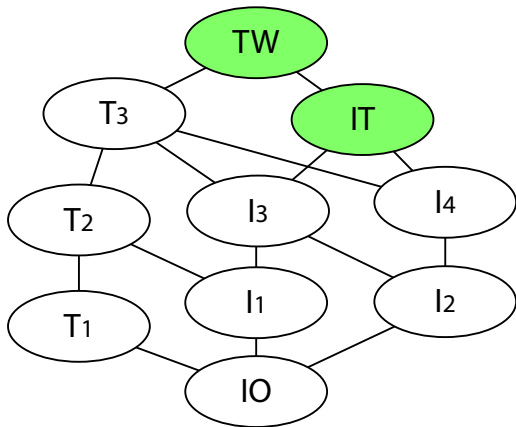
...in a way that is compatible with the simulated TW protocol.

Results: infinite memory



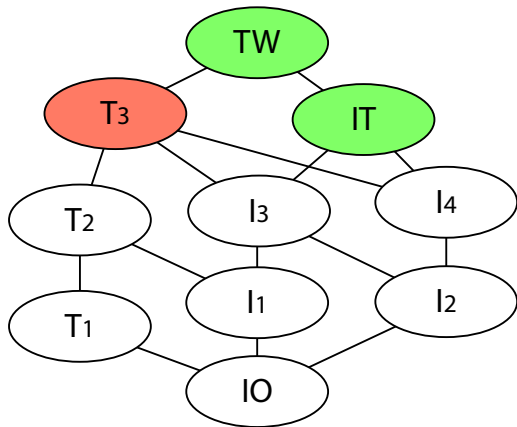
Suppose the simulating agents have **infinite memory**: what models can simulate *all* TW population protocols?

Results: infinite memory



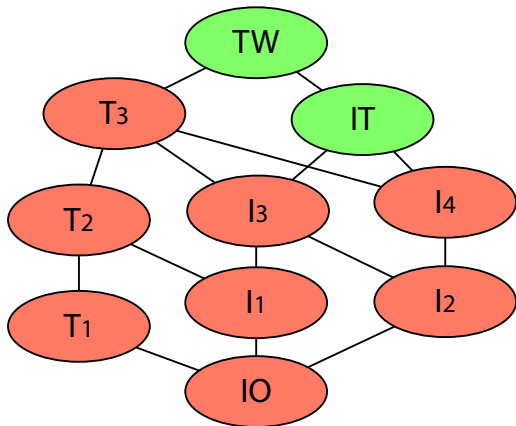
In **IT**, we can implement a *token-passing* technique that can be used to simulate two-way interactions.

Results: infinite memory



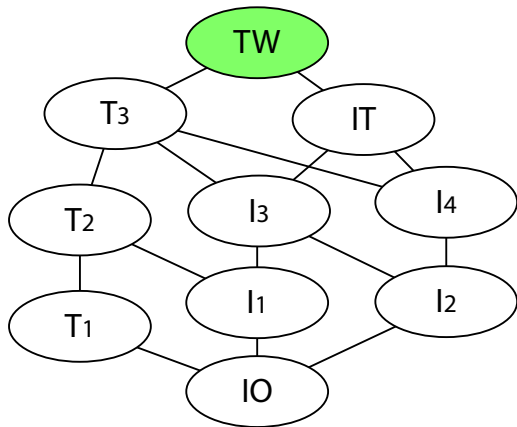
In T_3 , it is impossible to simulate a two-way protocol for the *pairing problem*.

Results: infinite memory



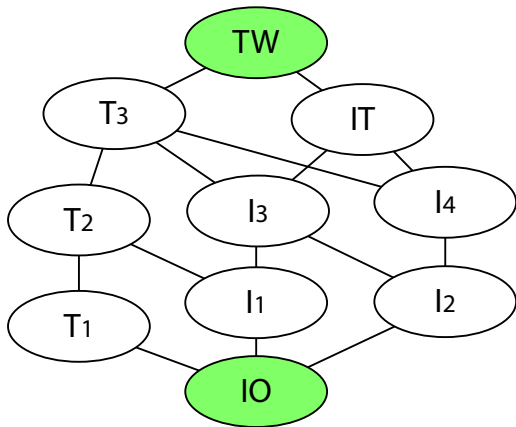
As a consequence, simulation is impossible also in the weaker interaction models.

Results: unique IDs



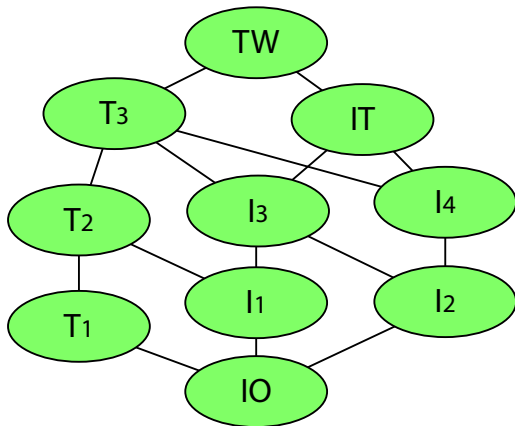
Suppose the simulating agents have *unique IDs* as part of their initial state.

Results: unique IDs



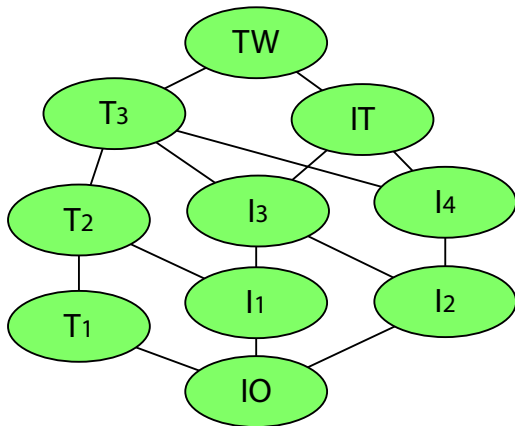
In **IO**, we can implement a *locking mechanism*, along with a *rollback process* to avoid deadlocks.

Results: unique IDs

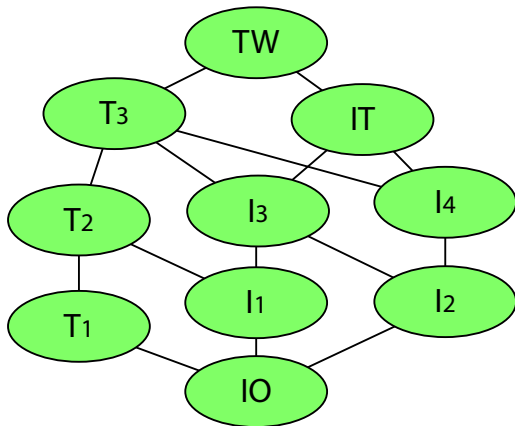


As a consequence, simulation is possible also in the stronger interaction models.

Results: knowledge of n

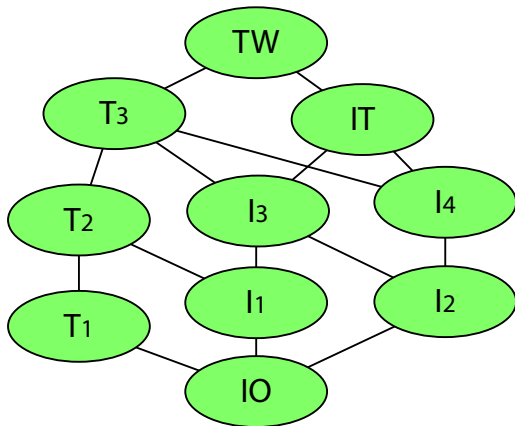


Suppose the simulating agents know the size of the system, n , and have $O(\log n)$ bits of internal memory.



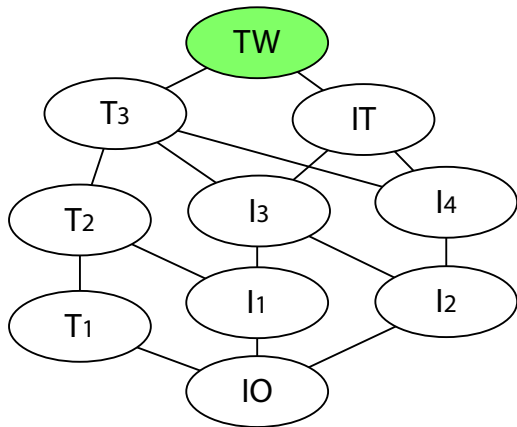
In **IO**, we can implement a *naming algorithm* that eventually gives each agent a unique ID.

Results: knowledge of n



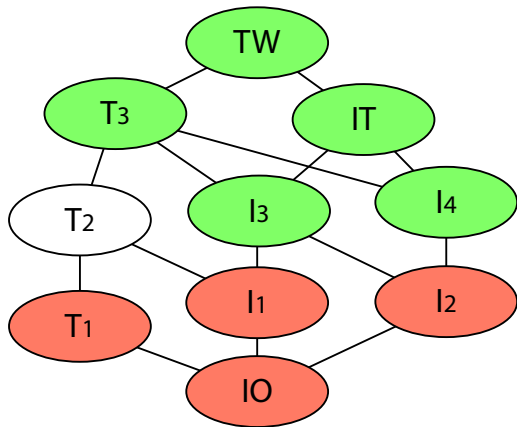
When an agent has ID n , the system starts executing the previous unique-ID simulation protocol.

Results: knowledge on omissions



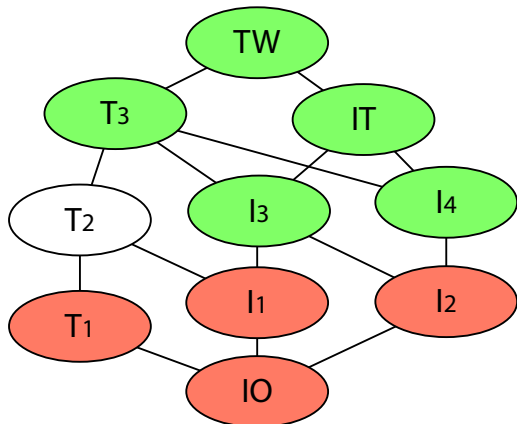
Suppose that the simulating agents are given an upper bound b on the number of faulty interactions in the system.

Results: knowledge on omissions



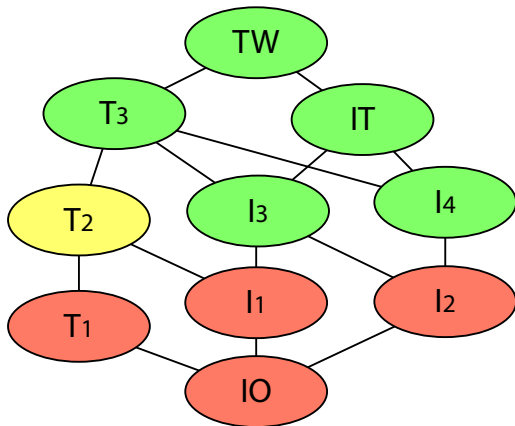
In I_3 and I_4 , we can extend the token-passing technique by splitting each token into $b + 1$ parts.

Results: knowledge on omissions



In \mathbf{T}_1 , \mathbf{I}_1 , and \mathbf{I}_2 , it is impossible to simulate the pairing protocol, even for $b = 1$.

Results: knowledge on omissions



Open problem: is it possible to simulate all TW protocols in T_2 , given an upper bound on the number of faulty interactions?